// YQIS - Paris INRIA

# Early Fault-Tolerant Quantum Algorithms in Practice

Oriel **Kiss**, November 6th, 2024

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Variational Algorithms (VQE, QAOA)

No theoretical arguments for quantum advantage

Scaling -  $\epsilon^{-2}$ 





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#### Fully-FTQC

Phase Estimation (QPE, QETU, filters, ...)

T-gate counts, High-depth Coherent computation

Scaling -  $\epsilon^{\text{-1}}$ 





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No theoretical arguments	Optimized Runtime, samples-	T-gate counts,
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Scaling - $\epsilon^{-2}$	qubits	Coherent computation
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- 3. Some degree of robustness against algorithmic errors.



**Problem 1.** Given a precision  $\delta > 0$  and lower bound on the overlap parameter  $\eta > 0$ , we seek to decide if

 $Tr[\rho\Pi_{\leq x-\delta}] < \eta \quad or \quad Tr[\rho\Pi_{\leq x+\delta}] > 0.$ (3)



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 $C(x) = p(x) * \Theta(x)$ 

convolution

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Fourier moments computed on the QC with a Hadamard test.

$$egin{aligned} G(x) &= rac{1}{2} + rac{2\mathcal{F}}{M} \sum_{i=1}^M \Big[ \operatorname{Re}[g_{k_i}( au)] \sin{(k_i x)} \ &+ \operatorname{Im}[g_{k_i}( au)] \cos{(k_i x)} \Big]. \end{aligned}$$







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 $\eta = \sum |\langle E_i | \Psi \rangle|^2,$ i < m

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Hamiltonian Simulation: system-specific; we use (asymptotical sub-optimal) product formula because:

- 1. no ancilla overhead
- 2. take advantage of locality
- 3. often much better than what is guaranteed
- 4. can scale better than qubitization in some regime





## // Our Workflow



X X N N D U

CDF smoothes out with the system size (random initial state)



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XXZ model



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Use ANOVA to test the statistical significance using F-test

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#### **// Key Insights II: Quantitative Resources**

For a given maximal runtime **D**  $O(\delta^{-1} \log \delta^{-1} \eta^{-1})^{1}$  and accuracy  $\epsilon$ , the number of samples **M** required to guarantee the correct result with probability  $1 - \vartheta$  is

$$\begin{split} M = & \left\lceil 2 \cdot \left[ \frac{2.07\pi^{-1}(\log 4D + 1) + 1}{\eta - 2\epsilon} \right]^2 \\ & \left[ \log \log \left( \frac{1}{\tau \epsilon} \right) + \log \left( \vartheta^{-1} \right) \right] \right\rceil, \end{split}$$



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To resolve bad initial states, we require more depth (and not only samples).



1. Wan, Berta and Campbell, Phys. Rev. Lett. 129,030503

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#### **// Numerical Simulations**

We consider a fully-connected Heisenberg model with random couplings over N spins

$$\mathcal{H} = \frac{1}{N} \sum_{i < j} \sum_{a \in \{x, y, z\}} J_a^{ij} \cdot \sigma_a^i \sigma_a^j, \text{ where } J_a^{ij} \sim \mathcal{N}(0, 1).$$



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For dynamics, we use a **second-order Trotter-Suzuki** with time step  $\Delta t = \tau/8$ . The circuit construction is based on SWAP networks.

We look at N=26 spins,

and use initial states prepared via DMRG (with low bond dimension).



#### // Key Insights III: DMRG(bd = 10) Initial State





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#### // Key Insights III: DMRG(bd = 10) Initial State





#### Depth matters most!

 $p_0 = 7x \ 10^{-6}$  (overlap)  $M = 10^{13}$  samples

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- 1. LT algorithm (and friends) bridges the gap between the NISQ and FTQC eras.
- 2. Instead of aiming at the true ground-state energy, one can concentrate on a finding the inflection point of the spectral CDF.
- 3. LT algorithms are able to improve on classical solutions, using only limited quantum resources (~10<sup>5</sup> samples).
- 4. If arbitrary precision is required, use quantum phase estimation. However, If a good approximation is enough, LT is a robust algorithm, which can be run in practice.







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### // Bonus: what can we do on NISQ?

- Blunt et al: variational dynamics + ZNE on a H<sub>3</sub> molecule (6 qub
- We focus on extracting Fourier moments of a nuclear EFT (4 qubits).

$$egin{aligned} egin{aligned} \left\langle \Psi_{0} 
ight| \hat{O}(ec{q})^{\dagger} e^{-iHj au} \hat{O}(ec{q}) \left| \Psi_{0} 
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angle \end{aligned}$$

• We use **purified echo verification** + various error suppresion techniques (twirling, dynamical decoupling, pulse efficient calibration, readout calibration).







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Unprepare the state and verify that it is the zero state



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Verification passed -> state contributes +/- 1 to the expectation value.

Otherwise: -> garbage



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#### **Purification of the ancilla**

$$ho = \lambda |\psi\rangle \langle \psi| + (1-\lambda) \sum_{k=2}^{2^N} p_k |\psi_k\rangle \langle \psi_k|.$$
Pure main component

Noisy components

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Noisy components

Without noise: the ancilla is pure after post-selection.

With noise: It is not. Extract the closest pure state from measurements.





Kiss et al , arXiv:2401.13048



x100 error reduction











## **Thank You!**

#### // Team





Utkarsh Azad



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David Wakeham



Juan Miguel Arrazola

arXiv: 2405.03754

arXiv:

2401.13048



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### Hadamard test



Initial state



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Hadamard test



Initial state

 $|\phi\rangle = \alpha |\bar{0}\rangle + \beta |\bar{0}^{\perp}\rangle.$ 



Hadamard test



Initial state

$$\begin{split} |\Phi\rangle &= \frac{1}{\sqrt{2}} \left( |\bar{0}\rangle \otimes |0\rangle + B^{\dagger} O_{k}^{\dagger} \mathcal{U}(j\tau) O_{l} B |\bar{0}\rangle \otimes |1\rangle \right) \\ &\equiv \frac{1}{\sqrt{2}} \left( |\bar{0}\rangle \otimes |0\rangle + |\phi\rangle \otimes |1\rangle \right), \end{split}$$

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$$\operatorname{Re}\{lpha\} = rac{\langle X_a 
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### // Sparse interpolation via compressive sensing

• Compressive sensing is an optimal technique to recover a signal from few measurements if we know a basis where the signal is sparse.

**Theorem 3.** Let U be the orthogonal conversion matrix between the measurement basis and the basis where the signal is sparse, with  $|U_{k,j}| \leq \mu(U)$ . If the number of samples m is chosen such that

$$m \ge C_0 \mu(U)^2 S \log d/\delta \tag{28}$$

$$m \ge C_0' \log^2 d/\delta,\tag{29}$$

for some constant  $C_0, C'_0$ , then every signal of sparsity S can be recovered with probability  $(1 - \delta)$ .

Assuming that we only have access to  $m \in \Omega$  samples, and that f is S-sparse in the  $\psi$  basis, the optimal solution is given by  $f^* = \Psi x^*$ , where  $x^*$  is the solution of the convex optimization problem

$$\min_{\tilde{x}\in\mathbb{R}^d} \|\tilde{x}\|_1 \text{ subject to } f_k = (\Psi\tilde{x})_k \,\forall k\in\Omega.$$
(31)

d = number of moments
S = number of non-zero components
(sparsity)



### // Numerics



- 1. Not good if used with importance sampling.
- 2. Does not work for extrapolation
- Always good if you need the whole signal (even with shot noise).



